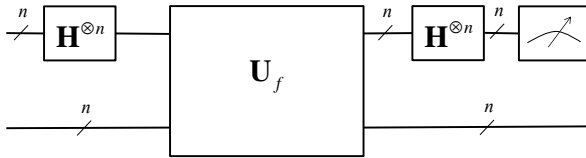


Simon's Periodicity Algorithm



PURPOSE: Detect the Period of a Function

1

Simon's Periodicity Algorithm Problem Overview

- Consider a Function of Interest of the Form:

$$f : \{0,1\}^n \rightarrow \{0,1\}^n$$

- Function is “Encoded” in an Oracle
- Compute a string $\mathbf{c}$,

$$\mathbf{c} = c_0 c_1 c_2 \dots c_{n-1}$$

Such that for all strings: $\mathbf{x}, \mathbf{y} \in \{0,1\}^n$

$$f(\mathbf{x}) = f(\mathbf{y}) \text{ if and only if } \mathbf{x} = \mathbf{y} \oplus \mathbf{c}$$

2

Simon's Periodicity Algorithm Problem Overview

$$f(\mathbf{x}) = f(\mathbf{y}) \text{ if and only if } \mathbf{x} = \mathbf{y} \oplus \mathbf{c}$$

- XOR ($\oplus$) operation is performed bitwise on the strings $\mathbf{y}$ and $\mathbf{c}$
- The values of f repeat themselves with respect to a Bitstring pattern $\mathbf{c}$
- $\mathbf{c}$ is the "Period" of f
- Purpose of Simon's Algorithm is to determine $\mathbf{c}$ (the "Period")

3

Reminder of the XOR Operator

x	y	$x \oplus y$
0	0	0
0	1	1
1	0	1
1	1	0

- $\oplus$, Denotes Addition Modulo-2:
 $(x + y)(\text{mod } 2) = x +_2 y = x \oplus y$
- $\oplus$, Referred to as the "Exclusive-OR" or XOR is Addition Modulo-2 When the Digit Set is Restricted to $\{0,1\}$
- XOR ($\oplus$) operation is performed bitwise on the n -bit strings $\mathbf{y}$ and $\mathbf{c}$
- EXAMPLES ($n = 3$):
 - $x = 011, y = 010: x \oplus y = 001$
 - $x = 101, y = 110: x \oplus y = 011$
 - $x = 111, y = 010: x \oplus y = 101$ *x inverts y*
 - $x = 000, y = 010: x \oplus y = 010$ *x creates identity for y*

4

Periodic Function Example

- number of bits $n=3$
- Consider $\mathbf{c}=101$

$$\overset{\mathbf{y}}{000} \oplus \overset{\mathbf{c}}{101} = 101 \Rightarrow f(000) = f(101)$$

$$001 \oplus 101 = 100 \Rightarrow f(001) = f(100)$$

$$010 \oplus 101 = 111 \Rightarrow f(010) = f(111)$$

$$011 \oplus 101 = 110 \Rightarrow f(011) = f(110)$$

$$100 \oplus 101 = 001 \Rightarrow f(100) = f(001)$$

$$101 \oplus 101 = 000 \Rightarrow f(101) = f(000)$$

$$110 \oplus 101 = 011 \Rightarrow f(110) = f(011)$$

$$111 \oplus 101 = 010 \Rightarrow f(111) = f(010)$$

this must hold if
 $f(\mathbf{x})=f(\mathbf{y})$ over $\mathbf{c}$,
 the period of f

if $\mathbf{c}=0^n$, what does
 that imply about f

5

Periodic Function Example

- number of bits $n=3$
- Consider $\mathbf{c}=101$

$$\overset{\mathbf{y}}{000} \oplus \overset{\mathbf{c}}{101} = 101 \Rightarrow f(000) = f(101)$$

$$001 \oplus 101 = 100 \Rightarrow f(001) = f(100)$$

$$010 \oplus 101 = 111 \Rightarrow f(010) = f(111)$$

$$011 \oplus 101 = 110 \Rightarrow f(011) = f(110)$$

$$100 \oplus 101 = 001 \Rightarrow f(100) = f(001)$$

$$101 \oplus 101 = 000 \Rightarrow f(101) = f(000)$$

$$110 \oplus 101 = 011 \Rightarrow f(110) = f(011)$$

$$111 \oplus 101 = 010 \Rightarrow f(111) = f(010)$$

this must hold if
 $f(\mathbf{x})=f(\mathbf{y})$ over $\mathbf{c}$,
 the period of f

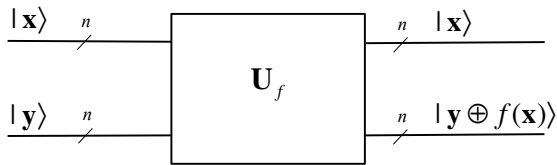
if only possible
 $\mathbf{c}=0^n$, what does
 that imply about f

f is not periodic

6

Function Represented as Oracle

- Function of Interest Specified as a Unitary Operation (Oracle) of the Form:



- Setting $y=0^n$ Provides a Convenient way to evaluate $f(x)$ When it is an Oracle
- Setting $y=1^n$ Provides a Convenient way to evaluate inverse of $f(x)$ When it is an Oracle

7

Classical Solution to Problem

- Evaluate $f(x)$ on different binary strings
- After Each Evaluation, Check if Function Response has Already Been Found
- If two strings x_1 and x_2 are found such that $f(x_1) = f(x_2)$ then it is assured that:

$$x_1 = x_2 \oplus c$$

- How do we find c ?

8

Classical Solution to Problem

- Evaluate $f(\mathbf{x})$ on different binary strings
 - After Each Evaluation, Check if Function Response has Already Been Found
 - If two strings $\mathbf{x}_1$ and $\mathbf{x}_2$ are found such that $f(\mathbf{x}_1) = f(\mathbf{x}_2)$ then: $\mathbf{x}_1 = \mathbf{x}_2 \oplus \mathbf{c}$
 - This Must Hold Consistently for All $\mathbf{x}_i$ Pairs
 - If it Does, How do we find $\mathbf{c}$?
- $$\mathbf{x}_2 \oplus \mathbf{x}_1 = \mathbf{x}_2 \oplus \mathbf{x}_2 \oplus \mathbf{c} = (\mathbf{x}_2 \oplus \mathbf{x}_2) \oplus \mathbf{c} = 0 \oplus \mathbf{c} = \mathbf{c}$$
- $$\mathbf{x}_1 \oplus \mathbf{x}_2 = \mathbf{c}$$

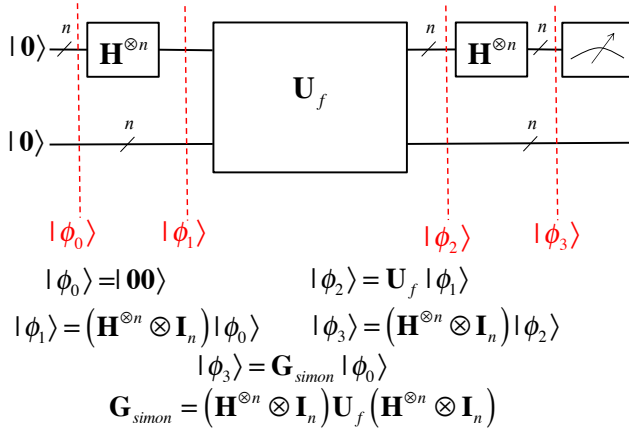
9

Classical Solution Complexity

- If $f(\mathbf{x})$ is periodic (it is two-to-one) then a repeat will be found before half of inputs are evaluated
 - If more than half of inputs checked with no match, then $f(\mathbf{x})$ is not periodic and $\mathbf{c} = 0^n$
 - Worst case number of evaluations:
- $$\frac{2^n}{2} + 1 = 2^{n-1} + 1$$
- Exponential Complexity

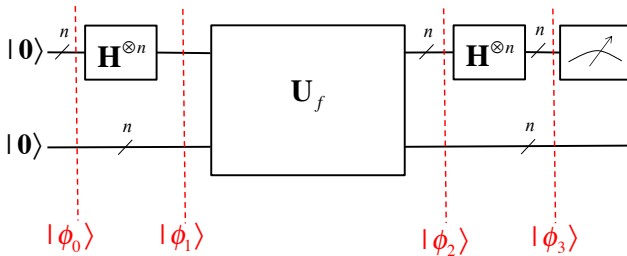
10

Periodicity Quantum Algorithm



11

Periodicity Quantum Algorithm

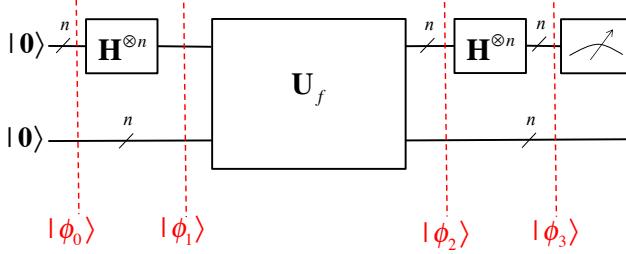


- Places Upper n Qubits in State of Superposition
- EXAMPLE: $n=3$:

$$|\phi_1\rangle = (H^{\otimes 3} \otimes I_3) |\phi_0\rangle = \frac{\sum_{\mathbf{x} \in \{0,1\}^3} |\mathbf{x}, 0\rangle}{\sqrt{2^3}} = \frac{|000,000\rangle + |001,000\rangle + \dots + |111,000\rangle}{\sqrt{8}}$$

12

Periodicity Quantum Algorithm



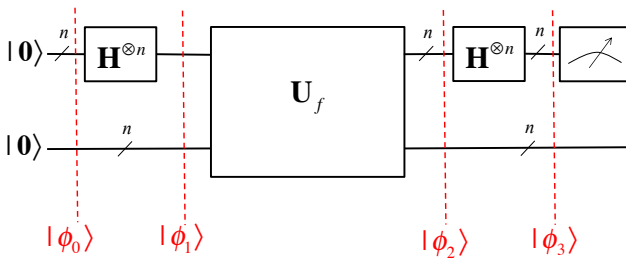
- U_f causes evaluation of f for the superimposed quantum state

$$|\phi_2\rangle = U_f(H^{\otimes n} \otimes I_n)|\phi_1\rangle = \frac{1}{\sqrt{2^n}} \sum_{x \in \{0,1\}^n} |x, f(x)\rangle$$

Hadamard is Applied to $|0\rangle$ so we Express it this way

13

Periodicity Quantum Algorithm



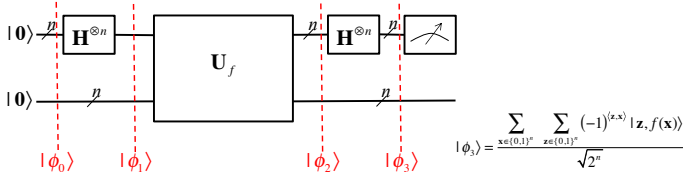
- n -Dimensional Hadamard Transform applied again

$$|\phi_3\rangle = (H^{\otimes n} \otimes I_n) |\phi_2\rangle = \frac{\sum_{x \in \{0,1\}^n} \sum_{z \in \{0,1\}^n} (-1)^{\langle z, x \rangle} |z, f(x)\rangle}{\sqrt{2^n}}$$

Hadamard is Expressed Generally as Explained in Class Notes on "Hadamard Transform"

14

Periodicity Quantum Algorithm



- n -Dimensional Hadamard Transform applied again
- For each x and z in the Summations, Consider when: $|z, f(x)\rangle = |z, f(x \oplus c)\rangle$
- When this occurs, same coefficients appear **twice in Double Summations**:

first instance for $f(x)$

second instance for $f(x \otimes c)$

same result from 2nd second set of Hadamard Gates after Oracle since $f(x) = f(x \otimes c)$

$$\begin{aligned} \frac{(-1)^{\langle z, x \rangle} + (-1)^{\langle z, x \oplus c \rangle}}{2} &= \frac{(-1)^{\langle z, x \rangle} + (-1)^{\langle z, x \rangle \oplus \langle z, c \rangle}}{2} \\ &= \frac{(-1)^{\langle z, x \rangle} + (-1)^{\langle z, x \rangle} (-1)^{\langle z, c \rangle}}{2} \end{aligned}$$

Make use of "Bit String Inner Product Properties"
Slide in Class Notes for "Hadamard Transform"
re: (.) Notation (refer to this slide set for reminder)

15

Periodicity Quantum Algorithm

$$\frac{(-1)^{\langle z, x \rangle} + (-1)^{\langle z, x \oplus c \rangle}}{2} = \frac{(-1)^{\langle z, x \rangle} + (-1)^{\langle z, x \rangle \oplus \langle z, c \rangle}}{2} = \frac{(-1)^{\langle z, x \rangle} + (-1)^{\langle z, x \rangle} (-1)^{\langle z, c \rangle}}{2}$$

- When: $\langle z, c \rangle = 1$ (Destructive Interference – Amplitudes go to 0)

$$\frac{(-1)^{\langle z, x \rangle} + (-1)^{\langle z, x \rangle} (-1)^{\langle z, c \rangle}}{2} = \frac{(-1)^{\langle z, x \rangle} + (-1)^{\langle z, x \rangle} (-1)^1}{2} = \frac{(-1)^{\langle z, x \rangle} - (-1)^{\langle z, x \rangle}}{2} = 0$$

- When: $\langle z, c \rangle = 0$ (Constructive Interference – Amplitude Magnitudes are Non-zero)

$$\frac{(-1)^{\langle z, x \rangle} + (-1)^{\langle z, x \rangle} (-1)^{\langle z, c \rangle}}{2} = \frac{(-1)^{\langle z, x \rangle} + (-1)^{\langle z, x \rangle} (-1)^0}{2} = \frac{(-1)^{\langle z, x \rangle} + (-1)^{\langle z, x \rangle}}{2} = \pm 1$$

- Therefore, when measuring the top qubits, we only find those binary strings where: $\langle z, c \rangle = 0$
- This Happens Because Zero-valued Hadamard Coefficients Cancel Out Cases for $\langle z, c \rangle$ (0-valued multiplier)
- We Want These Cases in Our Measurement Since It Allows Us to Solve for the Period (as you will see later)

16

Example Function

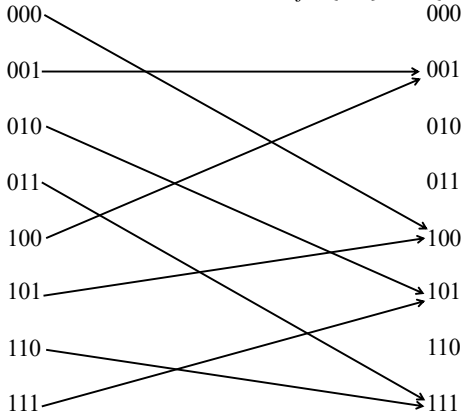
- Consider the Function: $f : \{0,1\}^3 \rightarrow \{0,1\}^3$
- Truth Table Representation – Embedded Inside U_f

$x_1 x_2 x_3$	f
0 0 0	100
0 0 1	001
0 1 0	101
0 1 1	111
1 0 0	001
1 0 1	100
1 1 0	111
1 1 1	101

17

Periodicity Algorithm Example

- Consider the Function: $f : \{0,1\}^3 \rightarrow \{0,1\}^3$



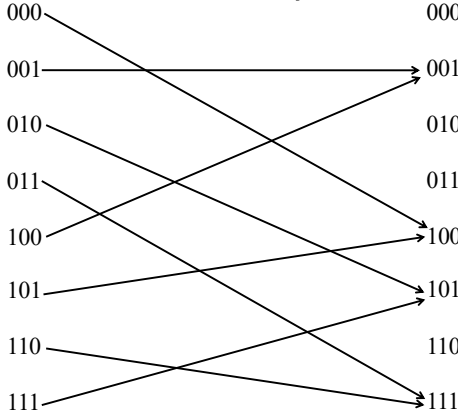
Is this Function
Periodic?

If "Yes," what is
the Period?

18

Periodicity Algorithm Example

- Consider the Function: $f : \{0,1\}^3 \rightarrow \{0,1\}^3$



Is this Function
Periodic?

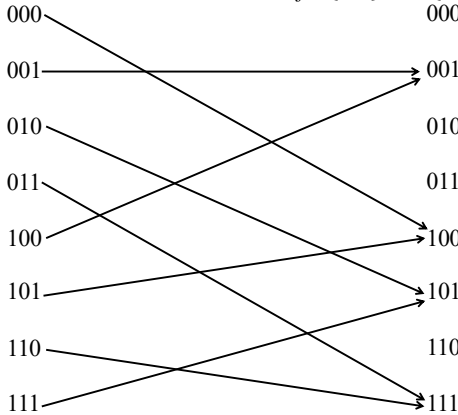
YES

If "Yes," what is
the Period?

19

Periodicity Algorithm Example

- Consider the Function: $f : \{0,1\}^3 \rightarrow \{0,1\}^3$



Is this Function
Periodic?

YES

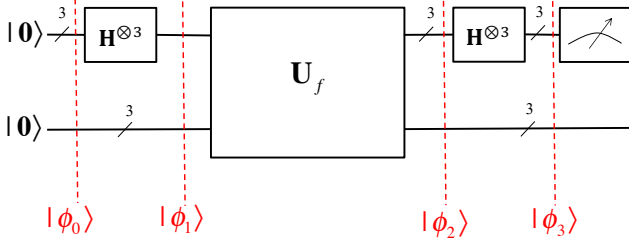
If "Yes," what is
the Period?

The Period is:

101

20

Periodicity Algorithm Example



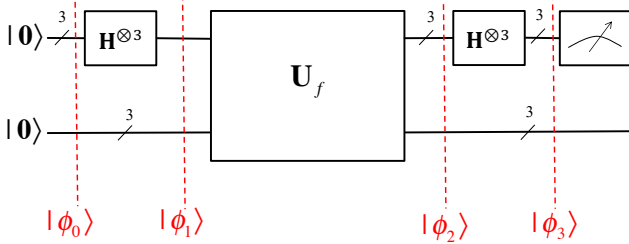
$$|\phi_0\rangle = |00\rangle = |0\rangle \otimes |0\rangle = |000\rangle \otimes |000\rangle$$

$$|\phi_1\rangle = (H^{\otimes 3} \otimes I_3) |\phi_0\rangle = \frac{\sum_{\mathbf{x} \in \{0,1\}^3} |\mathbf{x}, 0\rangle}{\sqrt{2^3}} = \frac{|000,000\rangle + |001,000\rangle + \dots + |111,000\rangle}{\sqrt{8}}$$

$$= \frac{1}{\sqrt{8}} (|000\rangle \otimes |000\rangle + |001\rangle \otimes |000\rangle + |010\rangle \otimes |000\rangle + |011\rangle \otimes |000\rangle + |100\rangle \otimes |000\rangle + |101\rangle \otimes |000\rangle + |110\rangle \otimes |000\rangle + |111\rangle \otimes |000\rangle)$$

21

Periodicity Algorithm Example



$$|\phi_2\rangle = U_f (H^{\otimes 3} \otimes I_3) |\phi_0\rangle = \frac{\sum_{\mathbf{x} \in \{0,1\}^3} |\mathbf{x}, f(\mathbf{x})\rangle}{\sqrt{2^3}} = \frac{\sum_{\mathbf{x} \in \{0,1\}^3} |\mathbf{x}\rangle \otimes |f(\mathbf{x})\rangle}{\sqrt{8}}$$

22

Periodicity Algorithm Example

$$|\phi_2\rangle = U_f(\mathbf{H}^{\otimes 3} \otimes \mathbf{I}_3)|\phi_0\rangle = \frac{\sum_{\mathbf{x} \in \{0,1\}^3} |\mathbf{x}, f(\mathbf{x})\rangle}{\sqrt{2^3}} = \frac{\sum_{\mathbf{x} \in \{0,1\}^3} |\mathbf{x}\rangle \otimes |f(\mathbf{x})\rangle}{\sqrt{8}}$$

$$|\phi_2\rangle = \frac{1}{\sqrt{8}} (|000\rangle \otimes |100\rangle + |001\rangle \otimes |001\rangle + |010\rangle \otimes |101\rangle + |011\rangle \otimes |111\rangle \\ + |100\rangle \otimes |001\rangle + |101\rangle \otimes |100\rangle + |110\rangle \otimes |111\rangle + |111\rangle \otimes |101\rangle)$$

• In the Next Stage of the Cascade:

$$|\phi_3\rangle = \frac{\sum_{\mathbf{x} \in \{0,1\}^3} \sum_{\mathbf{z} \in \{0,1\}^3} (-1)^{\langle \mathbf{z}, \mathbf{x} \rangle} |\mathbf{z}\rangle \otimes |f(\mathbf{x})\rangle}{(\sqrt{8})(\sqrt{8})}$$

23

Periodicity Algorithm Example

• Writing this out Term by Term Yields:

$$|\phi_3\rangle = \frac{1}{8} ((+1)|000\rangle \otimes |f(000)\rangle + (+1)|000\rangle \otimes |f(001)\rangle + (+1)|000\rangle \otimes |f(010)\rangle + (+1)|000\rangle \otimes |f(011)\rangle \\ + (+1)|000\rangle \otimes |f(100)\rangle + (+1)|000\rangle \otimes |f(101)\rangle + (+1)|000\rangle \otimes |f(110)\rangle + (+1)|000\rangle \otimes |f(111)\rangle \\ + ((+1)|001\rangle \otimes |f(000)\rangle + (-1)|001\rangle \otimes |f(001)\rangle + (+1)|001\rangle \otimes |f(010)\rangle + (-1)|001\rangle \otimes |f(011)\rangle \\ + (+1)|001\rangle \otimes |f(100)\rangle + (-1)|001\rangle \otimes |f(101)\rangle + (+1)|001\rangle \otimes |f(110)\rangle + (-1)|001\rangle \otimes |f(111)\rangle \\ + ((+1)|010\rangle \otimes |f(000)\rangle + (+1)|010\rangle \otimes |f(001)\rangle + (-1)|010\rangle \otimes |f(010)\rangle + (-1)|010\rangle \otimes |f(011)\rangle \\ + (+1)|010\rangle \otimes |f(100)\rangle + (+1)|010\rangle \otimes |f(101)\rangle + (-1)|010\rangle \otimes |f(110)\rangle + (-1)|010\rangle \otimes |f(111)\rangle \\ + ((+1)|011\rangle \otimes |f(000)\rangle + (-1)|011\rangle \otimes |f(001)\rangle + (-1)|011\rangle \otimes |f(010)\rangle + (+1)|011\rangle \otimes |f(011)\rangle \\ + (+1)|011\rangle \otimes |f(100)\rangle + (-1)|011\rangle \otimes |f(101)\rangle + (-1)|011\rangle \otimes |f(110)\rangle + (+1)|011\rangle \otimes |f(111)\rangle)$$

24

Periodicity Algorithm Example

$$\begin{aligned}
 &+ (+1)|100\rangle \otimes |f(000)\rangle + (+1)|100\rangle \otimes |f(001)\rangle + (+1)|100\rangle \otimes |f(010)\rangle + (+1)|100\rangle \otimes |f(011)\rangle \\
 &+ (-1)|100\rangle \otimes |f(100)\rangle + (-1)|100\rangle \otimes |f(101)\rangle + (-1)|100\rangle \otimes |f(110)\rangle + (-1)|100\rangle \otimes |f(111)\rangle \\
 \\
 &+ ((+1)|101\rangle \otimes |f(000)\rangle) + (-1)|101\rangle \otimes |f(001)\rangle + (+1)|101\rangle \otimes |f(010)\rangle + (-1)|101\rangle \otimes |f(011)\rangle \\
 &+ (-1)|101\rangle \otimes |f(100)\rangle + (+1)|101\rangle \otimes |f(101)\rangle + (-1)|101\rangle \otimes |f(110)\rangle + (+1)|101\rangle \otimes |f(111)\rangle \\
 \\
 &+ ((+1)|110\rangle \otimes |f(000)\rangle) + (+1)|110\rangle \otimes |f(001)\rangle + (-1)|110\rangle \otimes |f(010)\rangle + (-1)|110\rangle \otimes |f(011)\rangle \\
 &+ (-1)|110\rangle \otimes |f(100)\rangle + (-1)|110\rangle \otimes |f(101)\rangle + (+1)|110\rangle \otimes |f(110)\rangle + (+1)|110\rangle \otimes |f(111)\rangle \\
 \\
 &+ ((+1)|111\rangle \otimes |f(000)\rangle) + (-1)|111\rangle \otimes |f(001)\rangle + (-1)|111\rangle \otimes |f(010)\rangle + (+1)|111\rangle \otimes |f(011)\rangle \\
 &+ (-1)|111\rangle \otimes |f(100)\rangle + (+1)|111\rangle \otimes |f(101)\rangle + (+1)|111\rangle \otimes |f(110)\rangle + (-1)|111\rangle \otimes |f(111)\rangle
 \end{aligned}$$

- Coefficients in **Red** are the Elements of: $\mathbf{H}^{\otimes 3}$
- Evaluating the Function f in the Previous Equation Yields the Following:

25

Periodicity Algorithm Example

- Evaluating the Function f in the Previous Equation Yields the Following:

$$\begin{aligned}
 |\phi_3\rangle = \frac{1}{8} &((+1)|000\rangle \otimes |100\rangle + (+1)|000\rangle \otimes |001\rangle + (+1)|000\rangle \otimes |101\rangle + (+1)|000\rangle \otimes |111\rangle \\
 &+ (+1)|000\rangle \otimes |001\rangle + (+1)|000\rangle \otimes |100\rangle + (+1)|000\rangle \otimes |111\rangle + (+1)|000\rangle \otimes |101\rangle \\
 \\
 &+ ((+1)|001\rangle \otimes |100\rangle) + (-1)|001\rangle \otimes |001\rangle + (+1)|001\rangle \otimes |101\rangle + (-1)|001\rangle \otimes |111\rangle \\
 &+ (+1)|001\rangle \otimes |001\rangle + (-1)|001\rangle \otimes |100\rangle + (+1)|001\rangle \otimes |111\rangle + (-1)|001\rangle \otimes |101\rangle \\
 \\
 &+ ((+1)|010\rangle \otimes |100\rangle) + (+1)|010\rangle \otimes |001\rangle + (-1)|010\rangle \otimes |101\rangle + (-1)|010\rangle \otimes |111\rangle \\
 &+ (+1)|010\rangle \otimes |001\rangle + (+1)|010\rangle \otimes |100\rangle + (-1)|010\rangle \otimes |111\rangle + (-1)|010\rangle \otimes |101\rangle \\
 \\
 &+ ((+1)|011\rangle \otimes |100\rangle) + (-1)|011\rangle \otimes |001\rangle + (-1)|011\rangle \otimes |101\rangle + (+1)|011\rangle \otimes |111\rangle \\
 &+ (+1)|011\rangle \otimes |001\rangle + (-1)|011\rangle \otimes |100\rangle + (-1)|011\rangle \otimes |111\rangle + (+1)|011\rangle \otimes |101\rangle
 \end{aligned}$$

26

Periodicity Algorithm Example

- Evaluating the Function f in the Previous Equation Yields the Following (continued):

$$\begin{aligned}
 & + (+1)|100\rangle \otimes |100\rangle + (+1)|100\rangle \otimes |001\rangle + (+1)|100\rangle \otimes |101\rangle + (+1)|100\rangle \otimes |111\rangle \\
 & + (-1)|100\rangle \otimes |001\rangle + (-1)|100\rangle \otimes |100\rangle + (-1)|100\rangle \otimes |111\rangle + (-1)|100\rangle \otimes |101\rangle \\
 & + ((+1)|101\rangle \otimes |100\rangle + (-1)|101\rangle \otimes |001\rangle + (+1)|101\rangle \otimes |101\rangle + (-1)|101\rangle \otimes |111\rangle \\
 & + (-1)|101\rangle \otimes |001\rangle + (+1)|101\rangle \otimes |100\rangle + (-1)|101\rangle \otimes |111\rangle + (+1)|101\rangle \otimes |101\rangle \\
 & + ((+1)|110\rangle \otimes |100\rangle + (+1)|110\rangle \otimes |001\rangle + (-1)|110\rangle \otimes |101\rangle + (-1)|110\rangle \otimes |111\rangle \\
 & + (-1)|110\rangle \otimes |001\rangle + (-1)|110\rangle \otimes |100\rangle + (+1)|110\rangle \otimes |111\rangle + (+1)|110\rangle \otimes |101\rangle \\
 & + ((+1)|111\rangle \otimes |100\rangle + (-1)|111\rangle \otimes |001\rangle + (-1)|111\rangle \otimes |101\rangle + (+1)|111\rangle \otimes |111\rangle \\
 & + (-1)|111\rangle \otimes |001\rangle + (+1)|111\rangle \otimes |100\rangle + (+1)|111\rangle \otimes |111\rangle + (-1)|111\rangle \otimes |101\rangle)
 \end{aligned}$$

27

Periodicity Algorithm Example

- Combining Like Terms and Cancelling Out Where Possible Yields the Following:

$$\begin{aligned}
 |\phi_3\rangle &= \frac{1}{8} ((+2)|000\rangle \otimes |100\rangle + (+2)|000\rangle \otimes |001\rangle + (+2)|000\rangle \otimes |101\rangle + (+2)|000\rangle \otimes |111\rangle \\
 & + (+2)|010\rangle \otimes |100\rangle + (+2)|010\rangle \otimes |001\rangle + (-2)|010\rangle \otimes |101\rangle + (-2)|010\rangle \otimes |111\rangle \\
 & + (+2)|101\rangle \otimes |100\rangle + (-2)|101\rangle \otimes |001\rangle + (+2)|101\rangle \otimes |101\rangle + (-2)|101\rangle \otimes |111\rangle \\
 & + (+2)|111\rangle \otimes |100\rangle + (-2)|111\rangle \otimes |001\rangle + (-2)|111\rangle \otimes |101\rangle + (+2)|111\rangle \otimes |111\rangle) \\
 |\phi_3\rangle &= \frac{1}{8} ((+2)|000\rangle \otimes (|100\rangle + |001\rangle + |101\rangle + |111\rangle) \\
 & + (+2)|010\rangle \otimes (|100\rangle + |001\rangle - |101\rangle - |111\rangle) \\
 & + (+2)|101\rangle \otimes (|100\rangle - |001\rangle + |101\rangle - |111\rangle) \\
 & + (+2)|111\rangle \otimes (|100\rangle - |001\rangle - |101\rangle + |111\rangle))
 \end{aligned}$$

28

Periodicity Algorithm Example

$$|\phi_3\rangle = \frac{1}{8}((+2)|\mathbf{000}\rangle \otimes (|100\rangle + |001\rangle + |101\rangle + |111\rangle) \\ + (+2)|\mathbf{010}\rangle \otimes (|100\rangle + |001\rangle - |101\rangle - |111\rangle) \\ + (+2)|\mathbf{101}\rangle \otimes (|100\rangle - |001\rangle + |101\rangle - |111\rangle) \\ + (+2)|\mathbf{111}\rangle \otimes (|100\rangle - |001\rangle - |101\rangle + |111\rangle))$$

- Measuring the Top 3 Qubits Gives (with equal probability): $\mathbf{000}, \mathbf{010}, \mathbf{101}, \mathbf{111}$
- For Each of These Measured Quantum States, it is True that the Inner Product with the Period Bitstring $\mathbf{c}$ is Zero
- We want These Cases Only Because it Allows Us to Easily Solve for the Period $\mathbf{c}$ (*usually classically*)

29

Periodicity Algorithm Example

- For Each of These Measured Quantum States, it is True that the Inner Product with the Period Bitstring $\mathbf{c}$ is Zero
- The Algorithm/Circuit is Measured a Sufficient Number of Times to Ensure All Possible Measurements are Obtained
- This Yields a Set of Simultaneous Equations:

$$(i) \quad \langle 000, \mathbf{c} \rangle = 0$$

$$(ii) \quad \langle 010, \mathbf{c} \rangle = 0$$

$$(iii) \quad \langle 101, \mathbf{c} \rangle = 0$$

$$(iv) \quad \langle 111, \mathbf{c} \rangle = 0$$

30

Periodicity Algorithm Example

$$\langle 000, c_1 c_2 c_3 \rangle = 0$$

$$\langle 010, c_1 c_2 c_3 \rangle = 0$$

$$\langle 101, c_1 c_2 c_3 \rangle = 0$$

$$\langle 111, c_1 c_2 c_3 \rangle = 0$$

$$(0 \wedge c_1) \oplus (0 \wedge c_2) \oplus (0 \wedge c_3) = 0$$

$$(0 \wedge c_1) \oplus (1 \wedge c_2) \oplus (0 \wedge c_3) = 0$$

$$(1 \wedge c_1) \oplus (0 \wedge c_2) \oplus (1 \wedge c_3) = 0$$

$$(1 \wedge c_1) \oplus (1 \wedge c_2) \oplus (1 \wedge c_3) = 0$$

31

Periodicity Algorithm Example

$$c_2 = 0$$

$$c_1 \oplus c_3 = 0$$

$$c_1 \oplus c_2 \oplus c_3 = 0$$

- This Means that $c_1=c_3=0$ or that $c_1=c_3=1$
- We Know that $\mathbf{c}$ is NOT EQUAL to 000 Since Function was Found not to be One-to-One
- Therefore $c_1=c_3=1$
- Period of Function is:

$$\mathbf{c} = c_1 c_2 c_3 = 101$$

32

Periodicity Algorithm Example

- Must Run Simon's Algorithm Several Times to Measure n Different z Bitstrings
 - We Need a Complete Set of Linear Equations (modulo-2) in order to solve for each of the n distinct c_i bits in the "Period bitstring" c
- Generally Use a Classical Computer for Solving n Different Linear Equations (modulo-2)
- Simon's Algorithm Generally Viewed as a "Quantum Function" or "Quantum Subroutine"
- Could Use a Quantum Linear Solver Also but Must Consider Runtime and Cost

33

Simon's Alg. Intuition

- First Bank of Hadamard Gates Places All Function Domain Values Into Equal (Perfect) Superposition – Enables Oracle to Produce All Function Values in One Pass and in Superposition
- When Function is 2-to-1, Oracle Produces Two Identical and Distinct Superimposed Function Values for Each x : $f(x)$ and $f(x \oplus c)$
- Second Bank of Hadamard Gates Cause Constructive Interference for Two Identical and Distinct Function Values, so Probability Amplitudes Increase in Magnitude
 - and Destructive Interference for Function Pairs that Don't Match, so Probability Amplitudes Decrease in Magnitude (go to zero ideally)
- Several Runs Allows Different Function Values that Match to Result as Collapsed Values
 - We Keep Iterating to Obtain n Distinct $f(x_i)$ and $f(x_i \oplus c)$
 - Each Measurement Results in a Random Hyperplane Constraint Since it Collapses to a Single $f(x_i)$ and $f(x_i \oplus c)$
- These Collapsed Values Define Set of n Linear Equations that we can Solve for the Period
 - Since they are Orthogonal (zero-valued), we can solve a set of Linear Equations

34

Summary

- Quantum Computers are Good for:
 - 1) Searching Random Database (Grover's Algorithm)
 - 2) Finding Periodicity in Functions (Simon's & QFT)
 - 3) Factoring Semiprime Values (Shor's Algorithm)
 - 4) A Few Other Cases (finding QC algorithms is an active research area – but Hard to do!)
- Simon's Method is a Specific Example of Detecting Periodicity
 - Restricted to Functions of the Form:
$$f: \mathbb{B}^n \rightarrow \mathbb{B}^n$$
$$f: \{0,1\}^n \rightarrow \{0,1\}^n$$
- "Quantum Advantage" Known for Searching & Period-finding
- General Case for Period-finding is the "Quantum Fourier Transform" (QFT)
- Shor's Algorithm Depends on Efficient Period Finding of Special Function (the "Modulo-Powers Function")