## Hadamard Matrices/Operators



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# Hadamard Matrices

- Square Matrices with Mutually Orthonormal Rows/Columns
- All Matrix Elements are Either +1 or -1
- In Signal Processing, Known as the "Walsh Transform"
- Walsh Transform is Fourier Transform with Square Waves (Walsh Functions) as Basis Functions
  - Fourier Transform on Two-Element Additive Group  $\mathbb{Z}_2:(\{-1,+1\},+_2)$
- Different Row Orderings Yield Variations of the Walsh Matrix

# Hadamard Matrices

- Natural Row Ordering Defined by Outer/Tensor (Kronecker Product)
- Rademacher-Walsh Ordering Defined by XOR Operations among Adjacent Rows
- Transform can be Implemented Using *n*log*n* Operations ("Fast" Transform)

Can factor as sparse direct product factors

- Certain Forms can be Used Directly as Error Correcting Codes
- One Form is Known as the Reed-Muller Codes/Transform

#### Hadamard Matrix with Natural Ordering



- This Form uses Square Roots of Unity Shown as Points
   on the Unit Circle in the Complex Plane
- Transform is a Discrete Fourier Transform over GF(2)
- Can Think of this as a Discrete Fourier Transform with Discretized Orthogonal Square Wave Functions as the Basis Set

# **Fast Hadamard Transform**

 So-called "fast" transforms and Butterfly Diagrams (Signal Flow Graphs)



# Fast Hadamard Transform

So-called "fast" transforms due to Sparse





#### Hadamard Matrix Rademacher-Walsh Ordering



- This Form uses Boolean Logic Values Instead of Mappings to the Unit Circle in the Complex Plane
- Transform Yield a Form of ESOP in Classical Logic

# Rademacher-Walsh Transform

- Same as the Naturally-ordered Hadamard Transform with Rows/Columns Permuted
- Other Orderings Possible
- Referred to as "Walsh Transforms" in the Signal Processing Community
- Sometimes the Scale Factor  $(1/\sqrt{2})^n$  is Not Used in Signal Processing Applications

#### **Reed-Muller Matrix with Natural Ordering**



- This Form uses Boolean Logic Values Instead of Mappings to the Unit Circle in the Complex Plane
- Transform Yield a Form of ESOP in Classical Logic

#### **Reed-Muller Form of Hadamard Matrix**



 $e^{i\frac{2\pi(0)}{2}} = +1$  +1 Maps to Boolean 0 -1 Maps to Boolean 1

Use  $\mathbb{B}=\{0,1\}$  Instead of  $\mathbb{Z}_2$ 

- This Form Uses the Boolean Logic Values Instead of Mappings to the Unit Circle in the Complex Plane
- Transform Yields a Form of ESOP in Classical Logic

$$\mathbf{R}^{\otimes 3} = \mathbf{R}_3 = \mathbf{R}_1 \otimes \mathbf{R}_1 \otimes \mathbf{R}_1 = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \otimes \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \otimes \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \otimes \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}$$

# **Reed-Muller Form of Hadamard Matrix** $\mathbf{R}^{\otimes 3} = \mathbf{R}_{3} = \mathbf{R}_{1} \otimes \mathbf{R}_{1} \otimes \mathbf{R}_{1} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \otimes \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \otimes \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}$ $\mathbf{R}^{\otimes 3} = \bigotimes_{i=1}^{3} \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \otimes \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{bmatrix}$



For More Details of Classical Logic Synthesis of ESOPs: CSE 8387 Switching Theory Class

**Naturally Ordered Hadamard**  $\mathbf{H}^{\otimes 3} = \mathbf{H} \otimes \mathbf{H} \otimes \mathbf{H} = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \otimes \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \otimes \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$ Η Η Η

#### Hadamard & Superposition



- Alternative Notation (using symbolic logic)
  - Conjunctive Logic Operation (Binary AND function):  $\Lambda$

$$\mathbf{H} = \begin{bmatrix} h_{ij} \end{bmatrix} = \frac{1}{\sqrt{2}} (-1)^{i \wedge j}$$

- *i* and *j* are row and column numbers
- Can Rewrite Hadamard Matrix as:



 $\mathbf{H}^{\otimes 2} = \frac{1}{\sqrt{2}} \begin{bmatrix} (-1)^{0 \wedge 0} & (-1)^{0 \wedge 1} \\ (-1)^{1 \wedge 0} & (-1)^{1 \wedge 1} \end{bmatrix}_{1}^{0} \otimes \frac{1}{\sqrt{2}} \begin{bmatrix} (-1)^{0 \wedge 0} & (-1)^{0 \wedge 1} \\ (-1)^{1 \wedge 1} \end{bmatrix}_{1}^{0}$ 

• Multiply  $(-1)^x$  by  $(-1)^y = (-1)^{x+y}$ :

- Multiplying  $(-1)^x$  by  $(-1)^y = (-1)^{x+y}$ : Yields  $(-1)^{(2n)}$  or  $(-1)^{(2n)+1}$
- Where *n* is an Integer  $n \in \mathbb{Z}$ ,  $\mathbb{Z} = \{0, 1, 2, ...\}$
- Exponentiating -1 to a non-negative Integer Results in:

$$(-1)^{2n} = +1$$
  $(-1)^{2n+1} = -1$ 

• Therefore,  

$$(-1)^{(i \wedge j) + (k \wedge m)} = (-1)^{2n} = +1$$

$$(-1)^{(i \wedge j) + (k \wedge m)} = (-1)^{2n+1} = -1$$

- We Only Need to Consider if Exponent is Even or Odd
  - When:  $(i \wedge j) + (k \wedge m) = 2n \rightarrow (-1)^{(i \wedge j) + (k \wedge m)} = +1$
  - When:  $(i \land j) + (k \land m) = 2n + 1 \rightarrow (-1)^{(i \land j) + (k \land m)} = -1$

• This is Modulo-2 Addition, can Replace with Exclusive-OR

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- This is Modulo-2 Addition, can Replace with Exclusive-OR
  - When:  $(i \land j) + (k \land m) = 2n \longrightarrow (i \land j) \bigoplus (k \land m) = 0$
  - When:  $(i \land j) + (k \land m) = 2n + 1 \longrightarrow (i \land j) \bigoplus (k \land m) = 1$
- Can Express  $\mathbf{H}^{\otimes 2}$  Hadamard Matrix as:



# Notation

- Using this Form for Naturally Ordered Hadamard
  - Following Function Notation is Helpful
  - DO NOT Confuse with BraKet EXPECTED VALUE!!!!
  - The COMMA is Important (this is NOT  $\langle \mathbf{A} \rangle$  or  $\langle \Psi | \mathbf{A} | \Psi \rangle$ )  $\langle , \rangle : \{0,1\}^n \times \{0,1\}^n \rightarrow \{0,1\}$
  - Inner Product Function over Strings,  $\mathbb{B}^{2n} = \{0,1\}^{2n} \rightarrow \mathbb{B}$ **EXAMPLE:** Two *n*-bit Strings

 $\mathbf{x} = x_{n-1}x_{n-1}\cdots x_2x_1x_0$   $\mathbf{y} = y_{n-1}y_{n-1}\cdots y_2y_1y_0$ 

- Note, We Use Bit-wise Exclusive-OR (a string):  $\mathbf{x} \oplus \mathbf{y} = (x_{n-1} \oplus y_{n-1}), (x_{n-1} \oplus y_{n-1}), \dots, (x_1 \oplus y_1), (x_0 \oplus y_0)$
- The "Inner Product" Notation is a Single Value:  $\langle \mathbf{x}, \mathbf{y} \rangle = (x_{n-1} \land y_{n-1}) \oplus (x_{n-1} \land y_{n-1}) \oplus \cdots \oplus (x_1 \land y_1) \oplus (x_0 \land y_0)$

## **Bit-String Inner Product Properties**

$$\begin{array}{l} \langle \,,\,\rangle \colon \{0,1\}^n \times \{0,1\}^n \to \{0,1\} \\ \langle x,0\rangle = 0 & \langle 0,y\rangle = 0 \end{array} \end{array}$$

 $\langle \mathbf{x}, \mathbf{1} \rangle = x_{n-1} \oplus x_{n-2} \oplus \cdots \oplus x_1 \oplus x_0 \qquad \langle \mathbf{1}, \mathbf{y} \rangle = y_{n-1} \oplus y_{n-2} \oplus \cdots \oplus y_1 \oplus y_0$  $\langle \mathbf{x} \oplus \overline{\mathbf{x}}, \mathbf{y} \rangle = \langle \mathbf{x}, \mathbf{y} \rangle \oplus \langle \overline{\mathbf{x}}, \mathbf{y} \rangle \qquad \langle \mathbf{x}, \mathbf{y} \oplus \overline{\mathbf{y}} \rangle = \langle \mathbf{x}, \mathbf{y} \rangle \oplus \langle \mathbf{x}, \overline{\mathbf{y}} \rangle$  $\langle \mathbf{x} \oplus \overline{\mathbf{x}}, \mathbf{y} \rangle = \langle \mathbf{1}, \mathbf{y} \rangle \qquad \langle \mathbf{x}, \mathbf{y} \oplus \overline{\mathbf{y}} \rangle = \langle \mathbf{x}, \mathbf{1} \rangle$ 

$$\langle \mathbf{0} \wedge \mathbf{x}, \mathbf{y} \rangle = \langle 0^n, \mathbf{y} \rangle = 0 \quad \langle \mathbf{x}, \mathbf{0} \wedge \mathbf{y} \rangle = \langle \mathbf{x}, 0^n \rangle = 0 \langle \mathbf{1} \wedge \mathbf{x}, \mathbf{y} \rangle = \langle 1^n \wedge \mathbf{x}, \mathbf{y} \rangle = \langle \mathbf{x}, \mathbf{y} \rangle \quad \langle \mathbf{x}, \mathbf{1} \wedge \mathbf{y} \rangle = \langle \mathbf{x}, 1^n \wedge \mathbf{y} \rangle = \langle \mathbf{x}, \mathbf{y} \rangle$$

### Hadamard with New Notation

$$\mathbf{H}^{\otimes 2} = \left(\frac{1}{\sqrt{2}} \times \frac{1}{\sqrt{2}}\right) \begin{bmatrix} (-1)^{\langle 00,00 \rangle} & (-1)^{\langle 00,01 \rangle} & (-1)^{\langle 01,00 \rangle} & (-1)^{\langle 01,01 \rangle} \\ (-1)^{\langle 00,10 \rangle} & (-1)^{\langle 00,11 \rangle} & (-1)^{\langle 01,10 \rangle} & (-1)^{\langle 01,11 \rangle} \\ (-1)^{\langle 10,00 \rangle} & (-1)^{\langle 10,01 \rangle} & (-1)^{\langle 11,00 \rangle} & (-1)^{\langle 11,01 \rangle} \\ (-1)^{\langle 10,10 \rangle} & (-1)^{\langle 10,11 \rangle} & (-1)^{\langle 11,10 \rangle} & (-1)^{\langle 11,11 \rangle} \end{bmatrix} \begin{bmatrix} 0 \\ 0 \\ 0 \\ 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

• Now Can Write General Formula for:

$$\mathbf{H}^{\otimes n}(\mathbf{i},\mathbf{j}) = \frac{1}{\sqrt{2^n}} (-1)^{\langle \mathbf{i},\mathbf{j} \rangle}$$

- i and j are row and column numbers written as binary strings
- Quantum State Vector (example):

 $|\mathbf{0}\rangle = |000\cdots00\rangle = \begin{bmatrix} 1\\0\\0\\0\\0\\\vdots\\0\\0\\0\\111\cdots10\\111\cdots11 \end{bmatrix} \begin{array}{c} 000\cdots00\\000\cdots01\\000\cdots10\\\vdots\\111\cdots10\\111\cdots11 \end{array}$ 

Red strings are row numbers written as binary strings – NOT part of equation

• Multiplying a Quantum State Vector by  $\mathbf{H}^{\otimes n}$ 



- For Arbitrary Quantum State  $|\mathbf{y}\rangle$  $\mathbf{H}^{\otimes n}|\mathbf{y}\rangle = \mathbf{H}^{\otimes n}[-,\mathbf{y}] = \frac{1}{\sqrt{2^n}} \sum_{\mathbf{x} \in \{0,1\}^n} (-1)^{\langle \mathbf{x},\mathbf{y} \rangle} |\mathbf{x}\rangle$ 
  - Denotes "don't care" or All Possible Rows from 0 to *n*-1